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POLYNOMIAL MANIPULATION SYSTEM -
FORTRAN IV PROGRAM

PREPARED BY

THE COMPUTER CENTER

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FOREWORD

This report is a technical summary reporting the progress of a study conducted in the Mathematics Department and the Computer Center of Auburn University. The study is focused toward fulfillment of Contract No. DAAH01-68-C-0296 granted to Auburn University by the Army Missile Command, Huntsville, Alabama.

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ABSTRACT

A FORTRAN IV program which implements the Polynomial Manipulation System (PMS) is presented and described. PMS uses the Euclidean Algorithm to reduce a system of polynomials in several variables to a resultant system which can be solved sequentially as polynomials in one variable (Kronecker's method). PMS is described briefly and references are given to more complete discussions and to other pertinent literature.

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I. INTRODUCTION

The Polynomial Manipulation System (PMS) uses the Euclidean Algorithm for finding the eliminant and the greatest common divisor (g.c.d.) of two multi-variable polynomials. All polynomials involved are represented symbolically; PMS is a computer program whose input is the symbolic representation of two polynomials and whose output is (normally) the symbolic representation of their g.c.d. and eliminant.

The underlying theory and the application of PMS to the problem of solving systems of polynomial equations is discussed in [1], [2], and [3]. The present report describes a FORTRAN IV implementation of PMS developed on the IBM 360 Model 50 at Auburn University.

The program is described in Section II, the basic flow charts are given in Section III, Input/Output is discussed in Section IV, and efficiency of the method is discussed in Section V along with possible future work. The FORTRAN program is reproduced in Appendix A. Appendix B contains a simple example of the use of PMS for reduction of three polynomial equations in three variables to a resultant system which can be solved in sequence as polynomials in one variable.

II. PROGRAM DESCRIPTION

The PMS program is basically a main program with four subroutines, only one of which is significant. The other three subroutines are used for output, format headings on printed output and scaling of coefficients when they become large enough to possibly cause an overflow. In its present form the program is limited in that it is set up to use only 175K of IBM 360 storage. This limitation places constraints on the program which allows storage of only 50,000 polynomial entries (each term has $n + 1$ entries where n is the number of variables in the polynomial) which are presently set up as follows:

- 1) The pair of polynomials has, at most, four variables.
- 2) Each polynomial has at most 3160 terms.
- 3) The leading coefficient polynomials, to be defined below, can have, at most, 400 terms.

Minor modifications could increase the number of terms or variables or size of leading coefficients at the cost of decreasing the others or by use of a greater amount of machine storage. Still larger polynomials could be processed by use of tape, disc or other storage, but this has not been effected since such increases would only tend to accentuate certain disadvantages of the method to be discussed in Section V.

Consider the pair of polynomials $U, T : E^N \rightarrow R$. Let x_1, \dots, x_n denote the variables. Each of these polynomial functions can be considered as a polynomial in x_1 whose coefficients

would then be polynomial functions from E^{N-1} to R . Let U^0 and T^0 denote the polynomials in x_2, \dots, x_n which are the leading coefficients of U and T respectively considered as polynomials in x_1 , and let u and t denote the degrees of U and T in x_1 . We may assume $t \geq u$. Consider the polynomial R defined by

$$R = U^0 T - T^0 U x_1^{t-u}.$$

R is a polynomial in x_1, \dots, x_n . Considering R as a polynomial in x_1 with polynomial coefficients, it is seen that $\text{Degree}(R) < t$. Let $\text{Degree}(R) = r$. If $r \geq u$, let $T = R$ and repeat above procedure. If $r < u$, let $T = U$ and $U = R$ and repeat the above procedure. After a finite number of applications of this algorithm a polynomial R will be found whose degree in x_1 is zero. Thus R will be a polynomial in x_2, x_3, \dots, x_n . It is easily seen that at each stage R has any zeros that are common to U and T . The R which is free of x_1 is called the eliminant of U and T .

III. BASIC FLOW CHARTS

The flowcharts for output and scaling will be omitted as their detail is not significant to the main purpose of the program. The main program flow chart is given on page 4.

MAIN PROGRAM

Read U,T



Compute U^0, T^0



Print U,T



1



Determine which of U and T has greatest degree in x_1 . If it is T continue, if not interchange U and T, U^0 and T^0 .



Call RESIDUE
and Form R



Scale if necessary

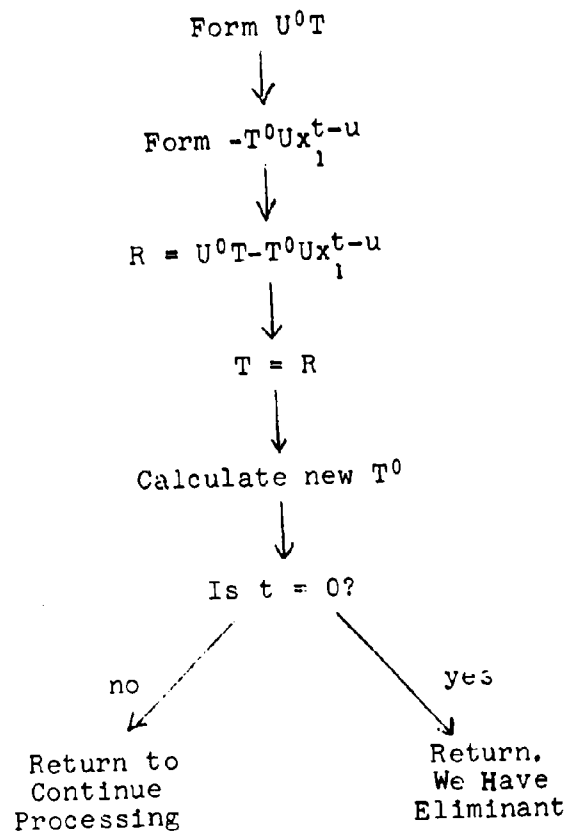


If R is free of x_1 , print R and return to beginning to read two new polynomials.
If not, let $T = R$.



1

RESIDUE SUBROUTINE



IV. INPUT/OUTPUT

The polynomials U and T are read in separately. The first card for each polynomial contains the number of variables in the polynomial in columns 31-34 in integer format, right justified. The number of terms in the polynomial appear in columns 35-38 in integer format, right justified. Following this card the terms of the polynomial appear, one term per card. The coefficient appears in columns 1-16 in E format; following this are the exponents of the variables right justified in integer format in columns 17-21, 22-26, 27-31, etc.

The output of this program is available in any medium, although the program is currently set up for printed output only.

V. EFFICIENCY AND FUTURE WORK

The PMS program has not proved useful as a method of reducing polynomial systems of equations to a resultant system for the following reasons:

- 1) Storage efficiency is low. An inordinate amount of core is needed to process many simple appearing systems of equations.

- 2) Time efficiency is low. Extreme amounts of time are needed to solve all but the most simple problems. As few as four equations in four variables with small exponents (on the order of ten or less) take many hours of machine time to reach a solution. Simpler problems are solvable in

small amounts of time, but other methods without these disadvantages can be used to solve these systems.

3) Certain types of systems give solutions which have a low order of accuracy. Several articles, [4], [5], [6], have been published discussing this problem as well as the two above.

The basic PMS program will be examined and modified to determine if it is of value in algebraically solving simple systems of differential equations.

VI. REFERENCES

1. Williams, L. H. Algebra of Polynomials in Several Variables for a Digital Computer, Duke University Memorandum 61-5-5,1 (1 May 1961).
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5. Collins, George T. Subresultants and Reduced Polynomial Remainder Sequences, J. ACM (Jan. 1967), 128-142.
6. Ku, S. Y. and Adler, R. J. Computing Polynomial Resultants: Bezout's Determinant vs. Collins' Reduced P.R.S. Algorithm, Comm. ACM (Jan. 1969), 23-30.


```

53      IUMAX(K-1,JU) = IR(K,J)
55      CONTINUE
        DO 10 J = 1, NTERM1
          U(J)=R(J)
          DO10I=1, MAX
10      IU(I,J)=IR(I,J)
          DO 9 J = 1, NTERM2
            R(J)=T(J)
            DO9I=1, NVAR2
9      IR(I,J)=IT(I,J)
          CALL PRINT(2, NTERM2, 0, 0, NVAR2)
          IF(NVAR2 .GE. NVAR1) GO TO 1075
          MAX = NVAR1
          NVAR2 = NVAR2 + 1
          DO 104 J = NVAR2, MAX
            DO 104 K = 1, NTERM2
104      IR(J,K) = 0
          NVAR2 = MAX
1075     MAX = NVAR2
          JPWRT = 0
          DO 60 J = 1, NTERM2
            IF(JPWRT - IR(1,J)) 51, 60, 50
51      JPWRT = IR(1,J)
60      CONTINUE
          JT = 0
          DO 56 J = 1, NTERM2
            IF(JPWRT - IR(1,J)) 56, 57, 56
57      JT = JT + 1
          TMAX(JT) = R(J)
          IF(ABS(TMAX(JT)).GE.1.E6) LM=1
          DO 58 K = 2, MAX
58      ITMAX(K-1, JT) = IR(K, J)
56      CONTINUE
          DO 8 J = 1, NTERM2
            T(J)=R(J)
            DO8I=1, MAX
8      IT(I,J)=IR(I,J)
107     IF(JPWRT-JPWRTU) 70, 71, 71
70      NN = JPWRT
          JPWRT = JPWRTU
          JPWRTU = NN
          MAXT=NTERM1
          IF(NTERM2.GT.NTERM1) MAXT=NTERM2
          DO80I=1, MAX
            TEMP=U(I)
            U(I)=T(I)
            T(I)=TEMP
            DO80J=1, MAXT
              TEMP=IU(J,I)
              IU(J,I)=IT(J,I)
80      IT(J,I)=TEMP
          NN = JU
          JU = JT
          JT = NN
          NN = NTERM1
          NTERM1 = NTERM2
          NTERM2 = NN
          NN = NVAR1

```

```

      NVAR1 = NVAR2
      NVAR2 = NN
      JJ=JU
      IF(JT.GT.JJ) JJ=JT
      MA = MAX-1
      DO 73 J = 1, JJ
      TW=UMAX(J)
      UMAX(J) = TMAX(J)
      TMAX(J)=TW
      DO73K=1,MA
      ITT=IUMAX(K,J)
      IUMAX(K,J) = ITMAX(K,J)
73    ITMAX(K,J)=ITT
71    CONTINUE
C
      IF(LM.EQ.1) CALL SCALE2(INTERM1,INTERM2,MAX)
      LM=0
101  CALL RESIDU (INTERM1,INTERM2,MAX,JPWRU,JPWRT,JU,JT)
      IF(LM.EQ.1) CALL SCALE2(INTERM1,INTERM2,MAX)
      IF(LM.EQ.1) GOTO 3001
      IF(LS2.EQ.1) CALL SCALE2(INTERM1,INTERM2,MAX)
      LS2=0
3001  LM=0
300  IF(LS1.EQ.0) GOTO 107
      LS1=0
      LM=0
301  CALL PRINT ( 3,JU,INTERM1,INTERM2,MAX)
      WRITE(6,10000) 1TIMES
10000  FORMAT(1H , 5HSCALE, 13)
      GO TO 100
1  FORMAT(30X,2I4)
4  FORMAT(E16.7,10I5)
      END

```

```

C      SUBROUTINE FOR PRINTING THE TWO POLYNOMIALS AND THE RESIDUE
C
C
C      SUBROUTINE PRINT (L,JU,NTERM1,NTERM2,N)
C*****
C*****
C
      IMPLICIT INTEGER*2(I-N)
      COMMON U,AX(400),IUMAX(3,400),ITMAX(3,400),TMAX(400),R(3160),
      NIR(4,3160),ITRATE,ITIMES,U(3160),IU(4,3160),T(3160)-IT(4,3160),
      NLS2,LS1,L4
C
C*****
C*****
C
      IF(L.EQ.3) WRITE(6,310)
      IF(L.EQ.2) WRITE(6,311)
      IF(L.EQ.1) WRITE(6,312)
50 CALL PRINTH(U)
      DO315K=1,JU
      WRITE(6,313) R(K),(IR(I,K),I=1,N)
315  CONTINUE
      K = N+1
503  CONTINUE
      RETURN
43  FORMAT(1X,F16.7,3X,10(15,5X))
311  FORMAT(1H1,'THE POLYNOMIAL U IS')
312  FORMAT(1H-,'THE POLYNOMIAL T IS')
310  FORMAT(1H-,'THE ELIMINANT IS')
540 FORMAT(10X,214)
541 FORMAT(16.7,1015)
      END

```

```

C SUBROUTINE FOR PRINTING HEADINGS
C SUBROUTINE PRINTH(K)

```

```

C THIS SUBROUTINE MERELY PRINTS COLUMN HEADINGS FOR THE
C VARIABLES DEPENDING ON THE NUMBER OF VARIABLES.
C THAT IS ITS ONLY PURPOSE. IT WILL HANDLE UP TO 10
C VARIABLES.

```

```

C GO TO (21,22,23,24,25,26,27,28,29,30),K
21 WRITE(6,31)
RETURN
22 WRITE(6,32)
RETURN
23 WRITE(6,33)
RETURN
24 WRITE(6,34)
RETURN
25 WRITE(6,35)
RETURN
26 WRITE(6,36)
RETURN
27 WRITE(6,37)
RETURN
28 WRITE(6,38)
RETURN
29 WRITE(6,39)
RETURN
30 WRITE(6,40)
RETURN
31 FORMAT(1H0,11HCOEFFICIENT,10X,4HX(1))
32 FORMAT(1H0,11HCOEFFICIENT,10X,4HX(1),5X,4HX(2))
33 FORMAT(1H0,11HCOEFFICIENT,10X,4HX(1),5X,4HX(2),5X,4HX(3))
34 FORMAT(1H0,11HCOEFFICIENT,10X,4HX(1),5X,4HX(2),5X,4HX(3),5X,4HX(4))
35 FORMAT(1H0,11HCOEFFICIENT,10X,4HX(1),5X,4HX(2),5X,4HX(3),5X,4HX(4),
1,5X,4HX(5))
36 FORMAT(1H0,11HCOEFFICIENT,10X,4HX(1),5X,4HX(2),5X,4HX(3),5X,4HX(4),
1,5X,4HX(5),5X,4HX(6))
37 FORMAT(1H0,11HCOEFFICIENT,10X,4HX(1),5X,4HX(2),5X,4HX(3),5X,4HX(4),
1,5X,4HX(5),5X,4HX(6),5X,4HX(7))
38 FORMAT(1H0,11HCOEFFICIENT,10X,4HX(1),5X,4HX(2),5X,4HX(3),5X,4HX(4),
1,5X,4HX(5),5X,4HX(6),5X,4HX(7),5X,4HX(8))
39 FORMAT(1H0,11HCOEFFICIENT,10X,4HX(1),5X,4HX(2),5X,4HX(3),5X,4HX(4),
1,5X,4HX(5),5X,4HX(6),5X,4HX(7),5X,4HX(8),5X,4HX(9))
40 FORMAT(1H0,11HCOEFFICIENT,10X,4HX(1),5X,4HX(2),5X,4HX(3),5X,4HX(4),
1,5X,4HX(5),5X,4HX(6),5X,4HX(7),5X,4HX(8),5X,4HX(9),5X,4HX(10))
END

```



```

C      RESIDUE SUBROUTINE      FORMS RESIDUE AND SORTS TERMS
C
      SUBROUTINE RESIDU(NTU,NTI,NV,JPU,JPT,JU,JT)
C*****
C*****
C
      IMPLICIT INTEGER*2(I-N)
      COMMON IUMAX(400),IUMAX(3,400),ITMAX(3,400),TMAX(400),R(3160),
      NIR(4,3160),ITRATH,ITIMES,U(3160),IU(4,3160),I(3160),IT(4,3160),
      NLS2,LS1,LM
      DIMENSION IB(5)
      MAXR=3160
      MAXP = 400
      INTEGER SWITCH
C
C*****
C*****
C
      I = 1
      DO 20 J = 1,NTI
      A=T(J)
      DO5210 K=1,NV
5210  IR(K)=IT(K,J)
      IF (IR(1).GE.JPT)GOTO20
      IF (ABS(A) - 1.E-6) 45,46,46
46  LS2=1
      RETURN
45  DO 10 L = 1, JU
      R(I) = IUMAX(L)*A
      IR(1,I) = IR(1)
      DO 10 KK = 2,NV
10  IR(KK,I) = IUMAX(KK-1,L) + IR(KK)
      MM=I-1
      IF (MM.EQ.0)GOTO10
      DO80KK=1,MM
      IF (IR(1,KK).NE.IR(1,I))GOTO80
      DO700KKK=2,NV
      IF (IR(KKK,KK).NE.IR(KKK,I))GOTO80
700  CONTINUE
      R(KK)=R(KK)+R(I)
      IF (ABS(R(KK)).GT.1.E-9)GOTO702
      NNM=MM-1
      DO900ILK=KK,NNM
      R(ILK)=R(ILK+1)
      DO800NPQ=1,NV
800  IR(NPQ,ILK)=IR(NPQ,ILK+1)
      I=I-1
702  I=I-1
      GOTO701
80  CONTINUE
701  IF (I.LT.MAXR)GOTO19
11  WRITE(6,12)
12  FORMAT(1H1,25HT00 MANY TERMS IN RESIDUE)
      STOP
19  I = I + 1
20  CONTINUE
      DO 40 J = 1,NTU
      A=U(J)

```

```

005211 K=1,NV
5211 IR(K)=IR(K,J)
IF(IR(1).GE.JPU)GOTO40
IF(ABS(A) - 1.E6) 47,48,48
48 LS2=1
RETURN
47 DO 41 L = 1, JT
R(I) = -TMAX(L)*A
IR(1,I) = IR(1) + JPT - JPU
DO 30 K = 2,NV
30 IR(K,I) = ITMAX(K-1,L) + IR(K)
MM=I-1
00350KK=1,MM
IF(IR(1,KK).NE.IR(1,I))GOTO380
003700KKK=2,NV
IF(IR(KKK,KK).NE.IR(KKK,I))GOTO380
3700 CONTINUE
R(KK)=R(KK)+R(I)
IF(ABS(R(KK)).GT.1.E-9)GOTO3702
NNM=MM-1
003800ILK=KK,NNM
R(ILK)=R(ILK+1)
003900NPD=1,NV
3800 IR(NPD,ILK)=IR(NPD,ILK+1)
I=I-1
I=I-1
3702 GOTO3701
380 CONTINUE
3701 IF(1.LT.4AXE)GOTO41
GOTO11
41 I = I + 1
40 CONTINUE
I = I - 1
99 JPWRT=0
JK=0
100 DO 105 M = 1,I
IF(IR(1,M) - JPWRT) 105,105,100
100 JPWRT = IR(1,M)
105 CONTINUE
JPT = JPWRT
IF(JPWRT)130,142,130
140 LS1=1
JL = I
RETURN
130 JT = 0
00711=1,400
TMAX(11)=0.
007JJ=1,3
ITMAX(JJ,11)=0
25 0056J=1,I
I(JPWRT - IR(1,J)) 56,57,56
57 JL = JT + 1
TMAX(JT) = I(J)
IF(TMAX(JT).GE.1.E6) LM=1
0053 K=2,NV
ITMAX(K-1,JT) = IR(K,J)
53 CONTINUE
56 CONTINUE

```

NOT REPRODUCIBLE

```

      DO 1 J=1,I
      T(J)=R(J)
      DO 1 L=1,NV
      IT(L,J)=IR(L,J)
1      NTT=I
376    RETURN
      END

```

```

      SUBROUTINE SCALE2(NTERM1,NTERM2,NV)
      IMPLICIT INTEGER*2(I-N)
      COMMON UMAX(400),IUMAX(3,400),ITMAX(3,400),TMAX(400),R(3160),
      NIR(4,3160),ITRATE,ITIMES,U(3160),IU(4,3160),T(3160),IT(4,3160),
      NLS2,LS1,LM
      EQUIVALENCE(NVAR1,MAX)
      MAX = NV
      NVAR2=NVAR1
      DO 1 I = 1,400
      UMAX(I) = UMAX(I)/1000.
1      TMAX(I) = TMAX(I)/1000.
      DO 7 J = 1,NTERM1
7      U(J) = U(J)/1000.
      DO 9 J = 1,NTERM2
9      T(J) = T(J)/1000.
      ITIMES = ITIMES + 1
      RETURN
      END

```

APPENDIX B

This appendix presents an example problem. The following system of three polynomials in three variables is reduced to a single polynomial in one variable:

$$(1) \quad x_1 + x_2 + x_3 = 0$$

$$(2) \quad x_1 + x_2^2 = 0$$

$$(3) \quad x_1 + x_3^2 = 0$$

First, x_1 is eliminated between (1) and (2) producing the following printout which has been labeled for expository convenience:

COEFFICIENT		x(1)	x(2)	x(3)	
0.1000000E	01	1	0	0	(1)
0.1000000E	01	0	1	0	
0.1000000E	01	0	0	1	
COEFFICIENT		x(1)	x(2)	x(3)	
0.1000000E	01	1	0	0	(2)
0.1000000E	01	0	2	0	
COEFFICIENT		x(1)	x(2)	x(3)	
0.1000000E	01	0	2	0	(E ₁)
-0.1000000E	01	0	1	0	
-0.1000000E	01	0	0	1	

Second, x_1 is eliminated between (3) and (2) producing the following printout:

COEFFICIENT		x(1)	x(2)	x(3)	
0.1000000E	01	1	0	0	(3)
0.1000000E	01	0	0	2	
COEFFICIENT		x(1)	x(2)	x(3)	
0.1000000E	01	1	0	0	(2)
0.1000000E	01	0	2	0	
COEFFICIENT		x(1)	x(2)	x(3)	
0.1000000E	01	0	2	0	(E ₂)
-0.1000000E	01	0	0	2	

Finally (E_1) and (E_2) are treated as a pair of polynomials in two variables x_1 and x_2 . Then x_1 (x_2 in our first system) is eliminated, producing the following printout:

COEFFICIENT		$x(1)$	$x(2)$	
0.1000000E 01	01	2	0	
-0.1000000E 01	01	1	0	(E_1)
-0.1000000E 01	01	0	1	
COEFFICIENT		$x(1)$	$x(2)$	
-0.1000000E 01	01	2	0	
0.1000000E 01	01	0	2	(E_2)
COEFFICIENT		$x(1)$	$x(2)$	
0.1000000E 01	01	0	4	
-0.2000000E 01	01	0	3	(E_3)

(E_3) is our eliminant free of x_1 and x_2 , so it can be solved. Using its solution (E_2) can then be solved. Using this (2) can be solved and the solutions to the resultant system (2), (E_2) , (E_3) are the solutions to the system (1), (2), (3). Three other equations from these six could have been taken to form the resultant system.

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APPENDIX C

Security Classification

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<p>A FORTRAN IV program which implements the Polynomial Manipulation System (PMS) is presented and described. PMS uses the Euclidean Algorithm to reduce a system of polynomials in several variables to a resultant system which can be solved sequentially as polynomials in one variable (Kronecker's method). PMS is described briefly and references are given to more complete discussions and to other pertinent literature.</p>		

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